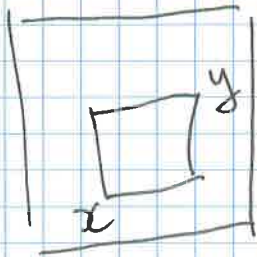


The largest empty box

~~The Problem~~

$$P_u \in [0, 1]^d \quad \#P_u = n$$



$$B = \{ [x, y] : x \leq y \}$$

$$\text{disp}(P_u) = \sup \{ |B| : B \cap P_u = \emptyset \}$$

Dispersion

$$\text{disp}(n, d) = \inf \{ \text{disp}(P_u) : \#P_u = n \}$$

Rok/Tichy 1996

$$n(\epsilon, d) = \min \{ n : \text{disp}(P_u) \leq \epsilon \}$$

$\epsilon > 0$

[Q] Study $\text{disp}(n, d)$, $n(\epsilon, d)$

$d \rightarrow \infty$

• MER

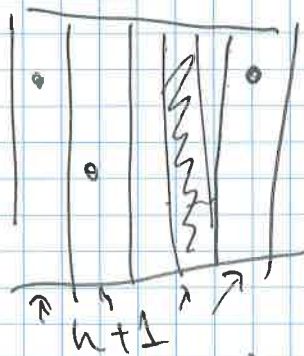
• Approximation of high-dim rank 1 tensors

2014 Bachmann

2015 Novak, Rudolf

- Plan
- ① nontrivial lower bound (Aistleitner, Rudolf)
 - ② ferhard Laderer: upper bound
 - ③ periodic disp. (Mario Allinetti)

① $disp(n, d) \geq \frac{1}{n+1}$ pigeon hole principle



$$\frac{1}{n+1}$$

$$n(\varepsilon, d) \geq \frac{1}{\varepsilon} - 1$$

• $disp(n, d) \geq \frac{5}{4(n+5)}$ Dimitrescu (Fian 2013)

$$n(\varepsilon, d) \geq \frac{5}{4\varepsilon} - 5$$

$$disp(4, d) \geq \frac{1}{4} + \text{pigeon hole. } d \geq 2$$

T#M: $disp(n, d) \geq \frac{\log_2 d}{4(n + \log_2 d)}$ $d \geq 32$

$$n(\varepsilon, d) \geq \frac{\log_2 d}{4\varepsilon} - \log_2 d$$

Proof: 1. Ingredient Pigeon Hole

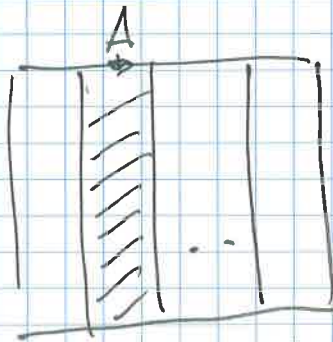
$$\text{disp}(u, d) \geq \frac{l+1}{n+l+1} \boxed{|\text{disp}(l, d)|} \quad (*)$$

$$\text{disp}(A, l, d) = \inf_{P_n \subseteq A} \sup |B| \quad \{B \subseteq A \cap P_n\}$$

$A \subseteq [0, 1]^d$ rect. box

$$\boxed{\text{disp}(A, l, d)} \geq |A| \text{disp}(l, d)$$

Pigeon hole: $n = (l+1)k + r \quad r \in \{0, 1, \dots, k-1\}$



$k+1$ boxes

one of these boxes (A) contains $\leq l$ points

$$\text{disp}(u, d) \geq \inf_{|A| = \frac{1}{k+1}} \text{disp}(A, l, d) \geq \frac{1}{k+1} \text{disp}(l, d)$$

$$k+1 = \frac{n-r}{l+1} + 1 \leq \frac{n+l+1}{l+1}$$

Impediment 2 non-trivial bound for (not too)

small l - structure in the coords.

(*) $disp(l, d) \geq \frac{1}{4}$

$2^l - 1 \leq d$

$P = \{t_1, \dots, t_l\} \subseteq \{0, 1\}^d$

$t_m \mapsto J_m \in \{0, 1\}^d$

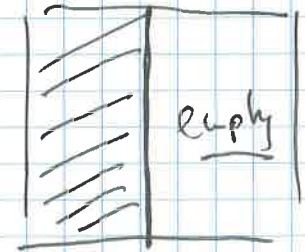
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$\frac{1}{2}$

Case 1: $J_m(i) = 0$ for some i

and all $m = \{1, \dots, l\}$

$disp \geq \frac{1}{2}$

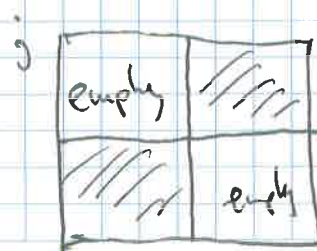


Case 2 $(J_m(i), \dots, J_m(i)) \in \{0, 1\}^l \setminus \{0, \dots, 0\}$

how many holes? $2^l - 2$

how many pigeons? $d \geq 2^l - 1$ $i = 1, \dots, d$

$\exists i \neq j \quad J_m(i) = J_m(j) \quad m = 1, \dots, l$



$disp \geq \frac{1}{4}$

$$\begin{aligned} \text{disp}(u, d) &\stackrel{(*)}{\geq} \frac{l+1}{u+l+1} \text{disp}(l, d) \\ l = \lfloor \log_2 d \rfloor &\stackrel{(**)}{\geq} \frac{1}{4} \frac{l+1}{u+l+1} \\ &= \frac{1}{4} \frac{\lfloor \log_2 d \rfloor + 1}{u + \lfloor \log_2 d \rfloor + 1} \\ &\geq \frac{1}{4} \frac{\log_2 d}{u + \log_2 d} \end{aligned}$$

$$\frac{d}{u}$$

② Upper bound (Gerhard Lander)

$$\text{disp}(u, d) \leq \frac{2}{u} \left(2^{cd} \right)$$

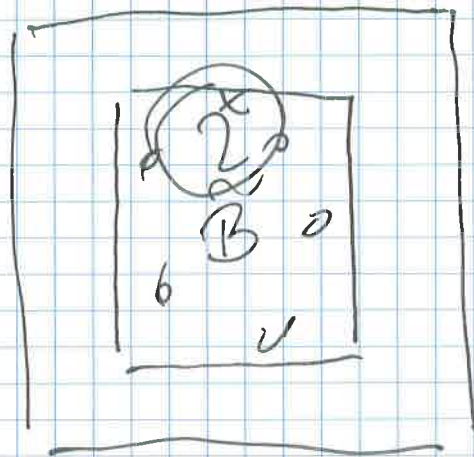
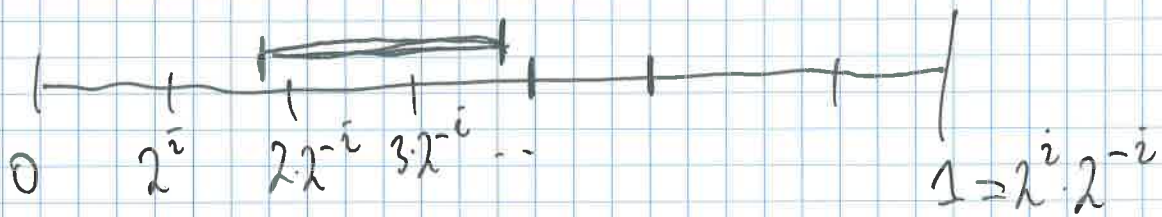
$$\frac{1}{u} \left(2^d \prod_{i=1}^{d-1} p_i \right)$$

Rok, Tidy 196

(t, u, d) - digital nets.

$b=2$, $t \leq cd$ (Niederreiter-Ying '96)

P_{2^m} : any elementary box of volume 2^{t-m}
contains exactly 2^t points.



$d=2 \quad |\tilde{B}| \geq \frac{1}{2} |B|$

\parallel
 2^{t-m}

$|B| \geq 2^{t-m} \cdot 2^d \Rightarrow |\tilde{B}| \geq 2^{t-m}$

$disp(P_{2^m}) \leq 2^{t-m+d} = \frac{1}{2^m} 2^{t+d}$

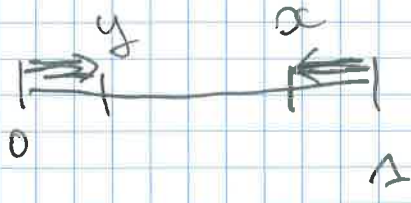
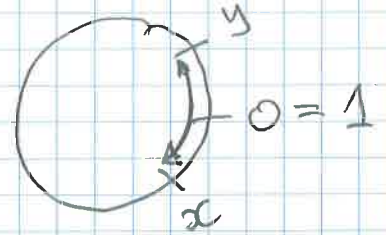
$2^m \leq n < 2^{m+1}$

$disp(P_n) \leq disp(P_{2^m}) \leq \frac{2}{n} 2^{6d}$

$\leq \frac{1}{2^m} 2^{(t+d)d}$
 $= \frac{1}{2^m} 2^{6d}$

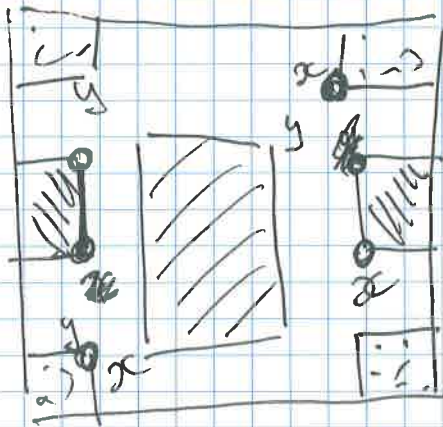
② Periodic Dispersion

$[0, 1] = \mathbb{T}^1$ torus



tensor products.

$B^{\mathbb{T}^d} = \{ [x, y] \mid x, y \in \mathbb{T}^d \}$



$disp^{\mathbb{T}^d}(P_n) = \sup_{P \cap B \neq \emptyset} \{ |B| : B \in B^{\mathbb{T}^d} \}$

$disp^{\mathbb{T}^d}(u, d), \quad n^{\mathbb{T}^d}(\epsilon, d)$

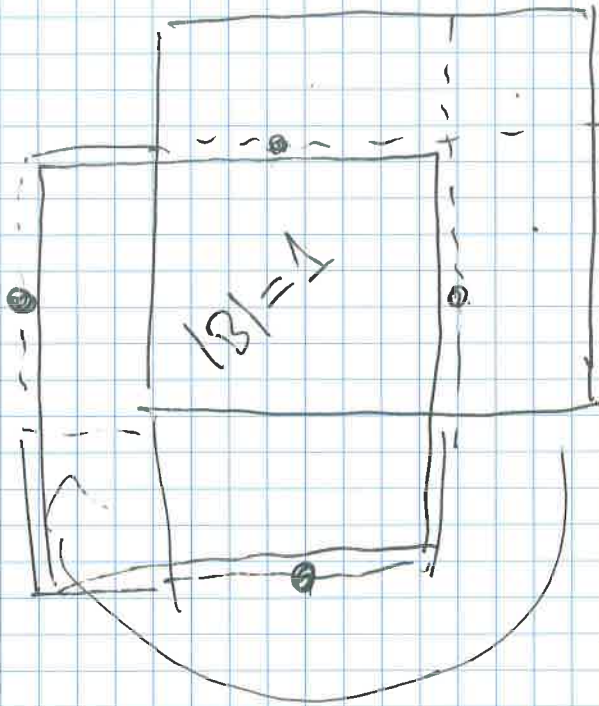
upper bound holds: $\leq \frac{2}{n} 2^{6d}$

$disp(u, d) \leq disp^{\mathbb{T}^d}(u, d)$

$disp^{\mathbb{T}^d}(u, d) \geq \min \{ 1, \frac{d}{n} \}$

$n^{\mathbb{T}^d}(\epsilon, d) \geq \frac{d}{\epsilon} \quad \epsilon < 1.$

Argument 1) ~~as~~ $n \leq d$, $n = d$

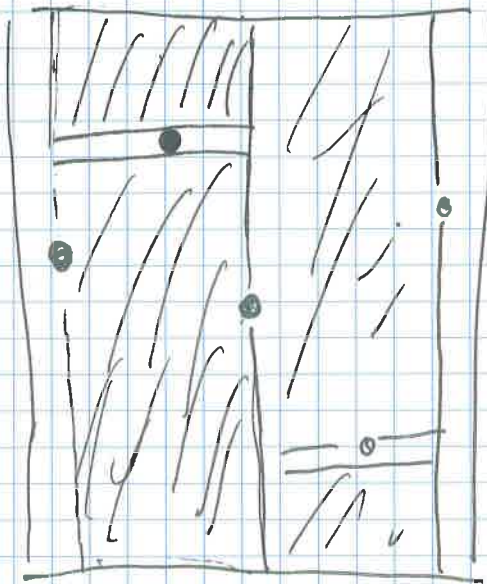


$$\underline{d = 2 = n}$$

$$|B| = 1$$

$$\text{disp}^n = 1 \text{ if } n \leq d.$$

$$n > d$$



$$t_1 \quad \dots \quad t_n$$

$$t_1(n) < t_2(n) < \dots < t_n(n)$$

$$\cdot 2 = \sum_{i=1}^n \text{empty boxes}$$

$$\text{One empty box} \geq \frac{2}{n} = d$$

2 Summary:

$$\frac{\log_2 d}{4(n + \log_2 d)} \leq \text{disp}(n, d) \leq \frac{2}{\epsilon} 2^{\epsilon d}$$

$$\text{with } \left\{ \frac{d}{n} \right\} \leq \text{disp}^{\#}(n, d) \leq \frac{1}{\epsilon} -$$

$$\frac{\log_2 d}{4\epsilon} - \log_2 d \leq n(\epsilon, d) \leq \frac{1}{\epsilon} \left(2^{6d+1} \right)$$

$$\frac{d}{\epsilon} \leq n^{\#}(\epsilon, d) \leq \frac{1}{\epsilon} - \quad \epsilon < 1$$

Compare this to the situation for L_{∞} -discrepancy.

$$\text{disc}^*(n, d) = \inf_{\#B_n = n} \|\Delta_{B_n}\|_{\infty}$$

$$n^*(\epsilon, d)$$

$$c_0 \text{ with } \left\{ \frac{d}{n} \right\} \leq \text{disc}_n^*(n, d) \leq 10 \sqrt{\frac{d}{n}}$$

$$\frac{d}{\epsilon} \leq n^*(\epsilon, d) \leq 100 \frac{d}{\epsilon^2} \quad \epsilon \leq \epsilon_0$$